

TAILORING THE LINUX KERNEL AND MPI TO A FAST AND NOISE-FREE MICROKERNEL

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HPC OS ARCHITECTURE



LWK

⊕ No Noise⊖ Compatibility

⊖ Features





Linux (tweaked)

⊙ Low Noise

- Compatibility
- ⊕ Features
- ⊖ Fast moving target



Argo



HPC OS ARCHITECTURE

AppAppLWK	LWK ⊕ No Noise ⊖ Compatibility ⊖ Features	CNK
App Proxy LWK Linux	 LWK + Linux ⊕ No Noise ⊕ Compatibility ⊕ Features ⊖ Effort 	mOS McKernel Hobbes/Kitten
AppAppLinux	 Linux (tweaked) Output Low Noise ⊕ Compatibility ⊕ Features ⊖ Fast moving target 	Cray Argo



VERSION 0 Pure L4





L4 Microkernel / Hypervisor





skampi - PingPong tests pt. I; 2 (of 10) procs



MPICH ON L4: PERFORMANCE



Tailoring the Linux Kernel and MPI to a Fast and Noise-Free Microkernel

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MPICH-L4| 0>--- p1MD version "0815" ---MPICH-L4| MPI_Wtime MPICH-L4| MPI_Wtime MPICH-L4| 0> [compiled: "2013-08-12"] MPICH-L4| 0> MPICH-L4 | 0><i> Simulation started :: 01.01.1970 00:00:03.993 (UTC +0) MPICH-L4| 0> MPICH-L4| 0> MPICH-L4| 0>Command: rom/p1md rom/md_10k_nochkpt.nml MPICH-L4| 0> MPICH-L4| 0> MPICH-L4 | Ø><i> 4 processes <more output here> MPICH-L4| 0><cpu> wall-clock timer function: MPI_Wtime MPICH-L4| 0> Work count sum: 184102828 MPICH-L4| 1> MPICH-L4| 2> MPICH-L4| 3><cpu> MD step performance: 2.936 steps/s MPICH-L4| 0> MD step performance: MPICH-L4| 1><cpu> 2.938 steps/s MPICH-L4| 2><cpu> MD step performance: 2.939 steps/s MPICH-L4| 3><cpu> MD interaction performance: 0.270 Mega interactions/s MD step performance: MPICH-L4| 0><cpu> 2.945 steps/s MPICH-L4| 1><cpu> MD interaction performance: 0.271 Mega interactions/s MPICH-L4| 2><cpu> MD interaction performance: 0.271 Mega interactions/s MPICH-L4| 0><cpu> MD interaction performance: 0.271 Mega interactions/s MPICH-L4 | 0> MPICH-L4| 0> MPICH-L4| 0><i> Simulation finished :: 01.01.1970 00:12:58.468 (UTC +0) MPICH-L4| 0>



- It worked:
 - p1MD, MPICH test, perftest, …
 - HPL, CP2K, ...
- Performance comparable to Linux^(*)
- Major limitation: no network backend



VERSION 1 Hybrid L4 + Linux











HYBRID SYSTEM



L4 Microkernel / Hypervisor



INFINIBAND DRIVER

Fast path (direct access to MMIO regs)

- Send / recv
- RDMA read / write
- Event polling

User Space

Kernel Space

Slow path (syscall required)

- Connection setup / teardown
- Buffer registration
- Event blocking





INFINIBAND DRIVER



L4 Microkernel / Hypervisor



VERSION 1: SPLIT MVAPICH2



Tailoring the Linux Kernel and MPI to a Fast and Noise-Free Microkernel

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VERSION 2Decoupled Threads



Linux

App

- Unmodified Linux programs (MPI, ...)
- L⁴Linux on L4 microkernel
- L4 microkernel controls the node
- Light-weight and low-noise if needed
- Linux process = L4 address space + thread
- Linux syscalls / exceptions: generic forwarding to L⁴Linux kernel





Linux

App

L⁴Linux

- Decoupling: move Linux thread to new L4 thread on its own core
- Linux syscall: Move back to Linux
- L4 syscalls:
 - Scheduling
 - Threads
 - Memory
- Direct I/O device access

L4 Microkernel / Hypervisor







Adam Lackorzynski, Carsten Weinhold, Hermann Härtig, "Decoupled: Low-Effort Noise-Free Execution on Commodity Systems", ROSS 2016, June 2016, Kyoto, Japan



DECOUPLING: BSP



Adam Lackorzynski, Carsten Weinhold, Hermann Härtig, "Decoupled: Low-Effort Noise-Free Execution on Commodity Systems", ROSS 2016, June 2016, Kyoto, Japan



VERSION 2.5 Decoupled Interrupts



BUSY WAITING





OVERSUBSCRIPTION





OVERSUBSCRIPTION

Unbalanced, no HT

Balanced (baseline), no HT









WAKE FROM IRQ: LINUX VS L4



Wake from interrupt on L4/Nova: 900 cycles, 0.3 µs

(best case, on Intel Core i7-4770 CPU @ 3.40GHz)

Adam Lackorzynski, Carsten Weinhold, Hermann Härtig, "Predictable Low-Latency Interrupt Response with General-Purpose Systems", OSPERT 2017, Duprovnik, Kroatia, June 2017



SUMMARY

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- L4 Microkernel + L⁴Linux + Decoupling
- Binary compatible + sysfs interface
- No modification to MVAPICH2
- Lessons learned:
 - Maximize reuse + minimize critical path
 - Methodology: Start from Linux, not from L4

