



High-Performance Big Data



DPU-Bench: A New Micro-Benchmark Suite to Measure the Offload Efficiency of SmartNICs

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(Paper: DPU-Bench: A Micro-Benchmark Suite to Measure Offload Efficiency Of SmartNICs

B. Michalowicz, K. Suresh, H. Subramoni, DK Panda, and S. Poole, Practice and Experience in Advanced Research Computing 23, Jul 2023)

Breakdown

- Introduction, Problem Statement, Motivation
- Design Choices
- Implementation
- Experimental Results
- Conclusion and Future Work

Problem Statement: We Need a New, DPU-Aware Microbenchmark Suite

- Already have plenty of Suites: OMB, IMB, mpiBench, OpenHPCA, etc.
 - NONE are DPU-Aware!
- Previous works: Utilize DPUs to offload in the context of an MPI-library, for specific apps, or in the context of Deep Learning, but no micro/benchmarks!
- SmartNICs are becoming popular: NVIDIA BlueField, AMD Pensando, Marvell etc.
- What algorithm is the "best" to offload to a DPU? Can we design a microbenchmark suite to measure the offload potential achieved from placing communication to the DPU?

Design Choice: Running at the IB-Verbs Level

- Running at the MPI level is bottlenecked by progress on the host
 - Naively offloading MPI-level ranks in, e.g., MPI_Ialltoall could create a bottleneck
 - Host issues MPI_Isends to a DPU process [] needs to check message progress from host AND DPU sides
- Running over IB-Verbs:
 - Closer to the hardware
 - DPU issues all RDMA operations host/CPU is not involved
 - No need for asynchronous progress on the host

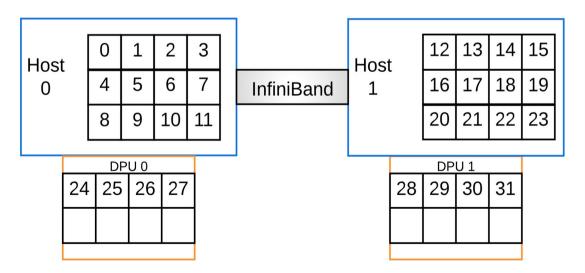
General Breakdown of DPU-Bench benchmarks (1/2)

- At initialization:
 - Exchange of metadata for RDMA Read/Write ops
 - Ikeys, rkeys, buffer addresses, QP numbers, etc.
- Runtime:
 - Step 1: Do Pure-Host execution [] obtain reference time for dummy compute.
 - Step 2: Offload communication to DPUs while the host side does compute
 - Step 3: Measure Offload Efficiency

Listing 2. General Approach to Each Benchmark	
/* Setup - assume options are	
passed in through CLI */	
<pre>MPI_Init();</pre>	Initialization
/*Record-keeping struct*/	
global_struct g;	
/* Every process makes its own m	emory region
* which makes the needed rkeys	
*/	
<pre>setup_ib_counters(&g, msg_size,</pre>	
<pre>num_workers, num_host_procs);</pre>	
<pre>/* Exchange of rkeys between processes */</pre>	
<pre>MPI_Allgather();</pre>	
/* Exchange of RDMA Buffer addresses */	
<pre>MPI Allgather();</pre>	
create_sends_and_recvs_pure_host	(); Step 1
<pre>create_sends_and_recvs_pure_host run_pure_host(); //gives referen</pre>	Step I
	ce time
<pre>run_pure_host(); //gives referen</pre>	ce time
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro</pre>	ce time
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()){</pre>	ce time cs();
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */</pre>	ce time cs();
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */ run_benchmark();</pre>	ce time cs(); Step 2
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */ run_benchmark(); }else{</pre>	ce time cs(); Step 2
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */ run_benchmark(); }else{</pre>	ce time cs(); Step 2
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<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */ run_benchmark(); }else{ perform_compute_on_host(ref) MPI_Barrier(MPI_COMM_WORLD);</pre>	Step 1 ce time cs(); Step 2 time);
<pre>run_pure_host(); //gives referen create_sends_and_recv_worker_pro if (proc_is_on_dpu()) { /* IB-level RDMA-operations */ run_benchmark(); }else{ perform_compute_on_host(ref_ } MPI_Barrier(MPI_COMM_WORLD); obtain_max_of_lat_and_comp();</pre>	Step 1 Step 2 time);

General Breakdown of DPU-Bench benchmarks (2/2)

- Assumptions made about MPI runtime
 - Use of a block-style hostfile helps with organization of config files
 - Higher-numbered ranks are placed on the DPU
 - Multiple Program, Multiple Data (MPMD) mode in an MPI library required for dealing with CPU/DPU configurations



Collective patterns in this benchmark suite

- Non-personalized one-to-all: Direct/Linear Broadcast
- Personalized all-to-one: Direct/Linear Gather
- Non-Personalized all-to-all: Direct/Linear Allgather with a single "root" worker
 - Later slides: Improve upon this to utilize multiple workers to demonstrate more efficient staging
- Will only show a subset here

Experiments Performed

- 8 nodes, 8 PPN on host
- Ranging from 1 worker (total) to 8 WPN on 8 DPUs (64 total workers)
 - Powers of 2: 1, 2, 4, 8...
- Study message sizes from 256K to 4 MB
- Offload Efficiency:

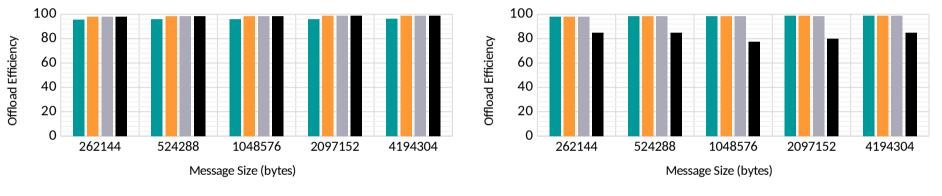
(reference_time/max(pure_comm, compute)) * 100

• Show Offload efficiency results for broadcast, gather, and allgather with both cyclic and block work

Offload Efficiency Results: Broadcast – Cyclic Work Assignment

Bcast Offload Efficiency (8 Nodes, 8 PPN) -- Cyclic Assignment

Bcast Offload Efficiency (8 Nodes, 8 PPN) --Cyclic Assignment





■ 2 WPN ■ 4 WPN ■ 6 WPN ■ 8 WPN

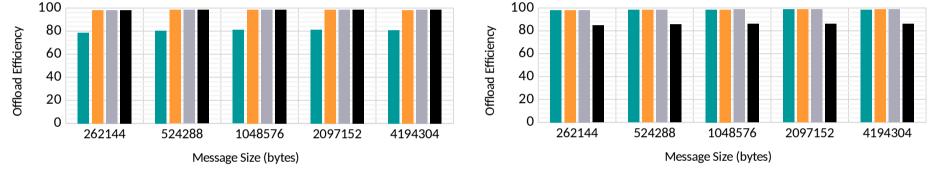
- General takeaway: 1 WPN is a "sweet spot" for maximum efficiency
- 8 WPN: Incurs overhead from the BF-2's single memory controller and limited cache sizes

Offload Efficiency Results: Gather - Cyclic Work

Assignment

Gather Offload Efficiency (8 Nodes, 8 PPN) --Cyclic Assignment

Gather Offload Efficiency (8 Nodes, 8 PPN) --Cyclic Assignment

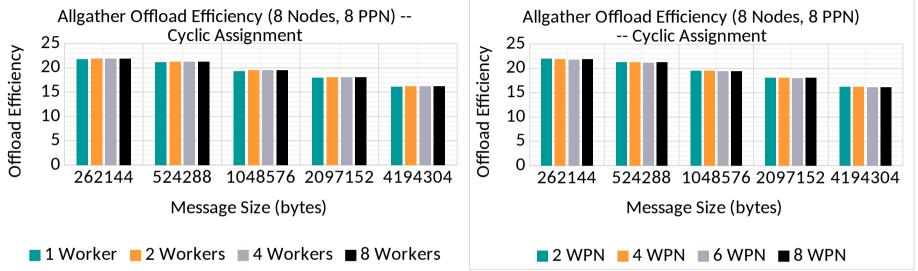


■ 1 Worker ■ 2 Workers ■ 4 Workers ■ 8 Workers

■ 2 WPN ■ 4 WPN ■ 6 WPN ■ 8 WPN

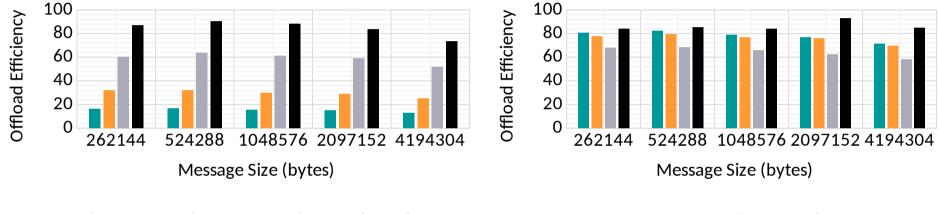
- 1 Worker: Adds overhead of an "intermediate" process to write to the root
- Similar trends to the "Broadcast" results

Offload Efficiency Results: Allgather (Single-Root Worker) – Cyclic Work Assignment



- Gather and Broadcast placed back-to-back incurs massive overhead
- Slight degradations with increase in message size

Offload Efficiency Results: Allgather (Efficient Use of Multiple Workers – Cyclic Work Assignment)



1 Worker 2 Workers 4 Workers 8 Workers

■ 2 WPN ■ 4 WPN ■ 6 WPN ■ 8 WPN

- 1 worker (total) to 1 WPN: 1.3-2X improvement in efficiency with the addition of workers
- 2 WPN 6 WPN: Predictable trend of decreasing offload efficiency

Conclusion/Future Work

- Introduction of a new, DPU-Aware Microbenchmark Suite
- Explored three communication patterns/algorithms
- Initial Release planned for the near future
- Generalize the benchmarks to other programming models
- Generalize the work to other SmartNICs

Thank You!

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Network-Based Computing Laboratory <u>http://nowlab.cse.ohio-state.edu/</u>



The High-Performance MPI/PGAS Project <u>http://mvapich.cse.ohio-state.edu/</u>



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